

# A Macro-Level Model for Investigating the Effect of Directional Bias on Network Coverage

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## Abstract

Random walks have been proposed as a simple method of efficiently searching, or disseminating information throughout, communication and sensor networks. In nature, animals (such as ants) tend to follow *correlated* random walks, i.e., random walks that are biased towards their current heading. In this paper, we investigate whether or not complementing random walks with directional bias can decrease the expected discovery and coverage times in networks. To do so, we develop a macro-level model of a directionally biased random walk based on Markov chains. By focussing on regular, connected networks, the model allows us to efficiently calculate expected coverage times for different network sizes and biases. Our analysis shows that directional bias can significantly reduce coverage time, but only when the bias is below a certain value which is dependent on the network size.

*Keywords:* Random walks, Markov chains, network coverage

## 1 Introduction

The concept of a random walk was introduced over a century ago by Pearson (Pearson 1905) and has been studied extensively since then (Dvoretzky & Erdős 1951, Brummelhuis & Hilhorst 1991, Caser & Hilhorst 1996). Recently, random walks have been proposed for searching, or disseminating information throughout, communications and sensor networks where the network's structure is dynamic, or for other reasons unknown (Bar-Yossef et al. 2006, Dolev et al. 2006, Avin & Brito 2004, Sadagopan et al. 2005). They are ideal for this purpose as they require no support information like routing tables at nodes (Avin & Krishnamachari 2006) — the concept of a random walk being for the agent performing the walk to move randomly to *any* connected node.

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The efficiency of random-walk-based algorithms can be measured in terms of the average number of steps the agent requires to cover every node in the network (and hence be guaranteed to find the target node in the case of search algorithms). This is referred to as the *coverage time* under the assumption that the agent takes one step per time unit. Obviously, improving the coverage time for algorithms is an important goal.

One straightforward approach to this is to have multiple agents (Alon et al. 2011). For some algorithms, such an approach is made even more effective when stigmergy is employed (Parunak 1997, Bonabeau et al. 1998, Ghosh et al. 2008), i.e., agents leave information for other agents directing them to their goals. Such an approach, inspired by the way ants leave trails of pheromones directing other ants to food (Camazine et al. 2001), is only useful when target nodes need to be visited by more than one agent. This is not always the case. More importantly, stigmergy is effective in directing agents only once a target node has been found. The time for the first agent to find a target node is not reduced. This can only be done by considering the agent's 'movement model'.

For this reason it has been suggested that random walks should be constrained, e.g., to prevent an agent returning to its last visited node, or to direct an agent to parts of the network where relatively few nodes have been visited (Lima & Barros 2007, Randall et al. 2007). We take a similar approach in this paper. We base our movement model on that observed in nature. Many models used by biologists to describe the movement of ants and other animals are based on *correlated* random walks, i.e., random walks which are biased to the animal's current direction (Kareiva & Shigesada 1983, Bovet & Benhamou 1988, McCulloch & Cain 1989). Based on our own observations of ants, we also investigate including a small probability of a non-biased step at any time to model occasional random direction changes.

To the best of our knowledge, the efficiency of directionally biased walks in networks have been investigated by only one other group of researchers. Fink et al. (Fink et al. 2012) look at the application of directional bias in a cyber-security system in which suspect malicious nodes must be visited by multiple agents. They compare coverage times for directional bias, with those for pure random walks, and random walks with stigmergy. Their conclusion is that directionally biased walks are more efficient even than random walks with stigmergy. This conclusion, however, is based on micro-level simulation, i.e., direct simulation of agents taking steps, for a single network size and bias. It cannot be generalised to arbitrary

network size or bias.

The micro-level simulation approach of Fink et al. requires coverage times to be calculated as the average of multiple runs. They performed 500 simulation runs for each movement model. Such an approach is impractical for a deeper investigation of the effect of directional bias which considers various network sizes and biases. For that reason, in this paper we develop a more abstract, macro-level model of a directionally biased walk. It builds on the work of Mian et al. (Mian et al. 2010) for random walks, describes the directionally biased walk in terms of a Markov chain (Norris 1998) and allows us to calculate the coverage time for a given network size and bias directly. Although certain special cases have analytic solutions, we have found this model to be helpful for a general approach to calculating coverage time.

This paper extends our previous work (Smith et al. 2014) by including a mathematical description of the macro-level model. We begin in Section 2 by describing the concept of a random walk, and our notion of directional bias in more detail. In Section 3 we present the Markov-chain model of a directionally biased walk on a network and show how it can be used to calculate the expected coverage time. In Section 4 we present and discuss the results of applying our model to the calculation of coverage times on a range of network sizes and biases. We conclude in Section 5.

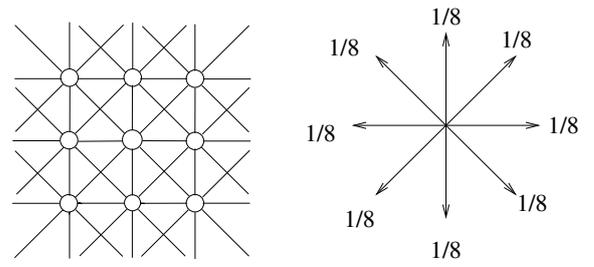
## 2 Directionally biased walks

Random walks have been studied in 1-dimensional, 2-dimensional and multi-dimensional spaces. Many of the results from 2-dimensional walks are applicable to communications and sensor networks which are commonly modelled as connected graphs. In particular, it is known that with probability 1 a random walk will cover every node of a connected graph (Aldous 1989) and a number of approaches for calculating the coverage time have been proposed (Aldous 1991, Dembo et al. 2004, Lima & Barros 2007, Mian et al. 2010).

The investigation in this paper focusses on regular, connected graphs where each node has exactly 8 neighbours (see Fig. 1(a) and, for the probability of the next step in a random walk over such a graph, Fig. 1(b)). Furthermore, to allow our graphs and hence networks to be finite, we wrap the north and south edges and the east and west edges to form a torus. Our aim is to provide a deeper analysis of directional bias than in the recent literature which also investigates the notion on regular toroidal graphs (Fink et al. 2012). While the results do not apply directly to arbitrary irregular networks, they do apply to random geometric graphs which are often used to model ad hoc sensor networks. Approaches using regular toroidal graphs to determine coverage time on random geometric graphs include that of Lima and Barros (Lima & Barros 2007) and Mian, Beraldi and Baldoni (Mian et al. 2010).

For modelling directional bias in nature, biologists typically use the von Mises distribution (Crist & MacMahon 1991). The von Mises distribution is a continuous angular function with a parameter  $\kappa$  which affects heading bias (see Figure 2).

We do not adopt the von Mises distribution in our approach for two reasons. Firstly, we have only a discrete number of directions and so do not require a continuous distribution. Secondly, as in random walks, we would like the computations the agent needs to perform to be simple. Our notion of directional bias limits our agent to choose either its current direction with a probability  $p$  (referred to as the *bias*),



(a) Regular, connected graph with 8 neighbours/node (b) Probability of next step on regular graph

Figure 1: Random walk on a regular, connected graph where each node has 8 neighbours.

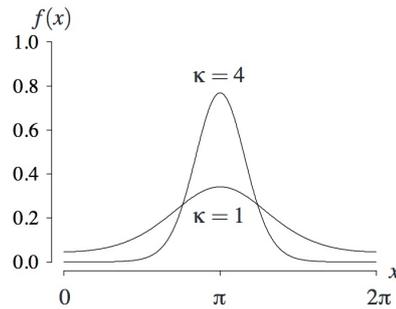


Figure 2: The von Mises distribution about  $\pi$  radians for  $\kappa = 1$  and  $\kappa = 4$ .

or any neighbouring direction, i.e.,  $\pi/4$  radians ( $45^\circ$ ) clockwise or anti-clockwise from the current direction, with equal probability of  $(1 - p)/2$ . When the bias  $p$  is high (as illustrated in Fig. 3(a)), the movement model approximates (discretely) that of the von Mises distribution for a high value of  $\kappa$  (such as  $\kappa = 4$  in Fig. 2). This is not the case, however, for low values of  $p$  (as illustrated in Fig. 3(b)).

Let  $Direction = \{0, \pi/4, \pi/2, 3\pi/4, \pi, 5\pi/4, 3\pi/2, 7\pi/4\}$  be the possible directions in radians. Our notion of directional bias is then defined formally as follows.

**Definition 1** (*Directional bias*) Given the current direction  $d \in Direction$  and bias  $p \in [0, 1]$ , the probability of moving in direction  $d' \in Direction$  at the next step,  $P(d')$ , is defined as follows.

$$\begin{aligned}
 d' = d &\Rightarrow P(d') = p \\
 d' \in \{(2\pi + (d - \pi/4)) \bmod 2\pi, (d + \pi/4) \bmod 2\pi\} &\Rightarrow P(d') = (1 - p)/2 \\
 d' \notin \{d, (2\pi + (d - \pi/4)) \bmod 2\pi, (d + \pi/4) \bmod 2\pi\} &\Rightarrow P(d') = 0 \quad \square
 \end{aligned}$$

We also investigate adding occasional random steps to our directionally biased walks. The idea is that with probability  $r$  the agent will make a random, rather than directionally biased, step. This better matches our own observations of the movements of ants.

**Definition 2** (*Directional bias with random steps*) Given the current direction  $d \in Direction$ , bias  $p \in [0, 1]$  and probability  $r \in [0, 1]$  of a random step, the probability of moving in direction  $d' \in Direction$  at the next step,  $P(d')$ , is defined as follows.

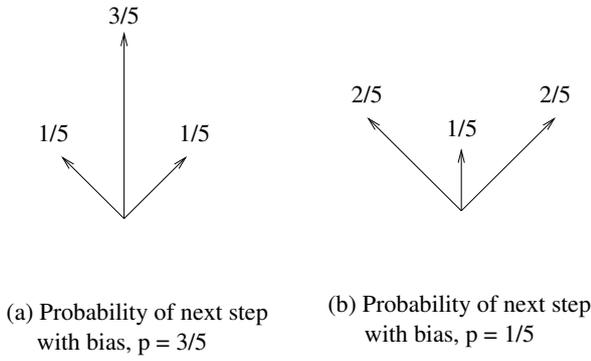


Figure 3: Directionally biased walk for agent with current direction north and biases  $p = 3/5$  and  $p = 1/5$ .

$$\begin{aligned}
 d' = d &\Rightarrow P(d') = r * 1/8 + (1 - r) * p \\
 d' \in \{(2\pi + (d - \pi/4)) \bmod 2\pi, (d + \pi/4) \bmod 2\pi\} \\
 &\Rightarrow P(d') = r * 1/8 + (1 - r)(1 - p)/2 \\
 d' \notin \{d, ((2\pi + (d - \pi/4)) \bmod 2\pi, (d + \pi/4) \bmod 2\pi)\} \\
 &\Rightarrow P(d') = r * 1/8 \quad \square
 \end{aligned}$$

To analyse coverage time under our models of directional bias, we adapt a Markov-chain model (Norris 1998) developed for random walks by Mian et al. (Mian et al. 2010). As explained in Section 3, this allows us to calculate the coverage time directly, and hence compare coverage times for different network sizes and biases.

### 3 A macro-level model

The previous work on directional bias by Fink et al. (Fink et al. 2012) shows that directionally biased walks are more efficient than random walks on a regular, connected toroidal graph; but only for the one specific network size and bias considered in their paper. In this paper, we produce more general results by investigating the effect on coverage time of varying the network size and directional bias. The micro-level model and simulation approach used by Fink et al. is not suited to this goal, requiring numerous simulations runs to calculate the coverage time for each network size and bias. We therefore use a more abstract, macro-level model which allows us to calculate the coverage time for a given network size and bias directly.

Our model is based on the work of Mian et al. (Mian et al. 2010) who provide a Markov-chain approach (Norris 1998) to model and calculate coverage time for random walks on a regular, connected toroidal graph. Given a network of  $N$  nodes, let the vector  $v$  of length  $N$  denote the state probability distribution with elements  $v_i$  for  $0 \leq i < N$ , and the matrix  $M$  of size  $N \times N$  denote the transition probability matrix with elements  $M_{i,j}$  for  $0 \leq i, j < N$ .

**Definition 3** (*Markov-chain model*) Let  $v^{(0)}$  denote the initial state distribution and  $v^{(k)}$  denote the state distribution after the  $k$ th step.

- The elements of the state probability distribution sum to 1.

$$\forall k \geq 0 \bullet \sum_{i=0}^{N-1} v_i^{(k)} = 1$$

- The rows of the transition probability matrix sum to 1.

$$\forall i < N \bullet \sum_{j=0}^{N-1} M_{i,j} = 1$$

- The state distribution at step  $k$  is calculated by multiplying the initial distribution by the transition probability matrix  $k$  times.

$$\forall k > 0 \bullet v^{(k)} = v^{(0)} M^k \quad \square$$

Often the system described by such a model would begin in a particular node with probability 1.

$$\exists i < N \bullet v_i^{(0)} = 1 \wedge (\forall j < N \bullet j \neq i \Rightarrow v_j^{(0)} = 0)$$

For a random walk, we call this node the *starting node*. A random walk is specified by letting the transition from a node to any of its neighbours occur with a probability of  $1/n$  where  $n$  is the number of neighbours. For example, for a 1-dimensional network of 5 nodes such that

$$v^{(0)} = (0 \quad 0 \quad 1 \quad 0 \quad 0)$$

the transition probability matrix for a random walk would be

$$M = \begin{pmatrix} 0 & \frac{1}{2} & 0 & 0 & \frac{1}{2} \\ \frac{1}{2} & 0 & \frac{1}{2} & 0 & 0 \\ 0 & \frac{1}{2} & 0 & \frac{1}{2} & 0 \\ 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} \\ \frac{1}{2} & 0 & 0 & \frac{1}{2} & 0 \end{pmatrix}$$

where, for example, row 0 (the topmost row of  $M$ ) indicates that an agent at node 0 (the first node of  $v$ ) has a probability of  $1/2$  of moving to node 1 and a probability of  $1/2$  of moving to node 4 (since we wrap the east and west edges). Multiplying  $v^{(0)}$  by  $M$  results in

$$\begin{aligned}
 v^{(1)} &= (0 \quad \frac{1}{2} \quad 0 \quad \frac{1}{2} \quad 0) \\
 v^{(2)} &= (\frac{1}{4} \quad 0 \quad \frac{1}{2} \quad 0 \quad \frac{1}{4})
 \end{aligned}$$

and so on. For a 2-dimensional network of  $n \times m$  nodes, the vector  $v$  would have  $n * m$  elements, with those from  $i * m \dots i * m + m - 1$  for  $0 \leq i < n$  denoting the nodes in the  $i$ th row of the network. So, for example, row 0 of the matrix for a random walk on a network of  $5 \times 5$  nodes is shown in Figure 4.

In order to calculate coverage time, we modify the standard Markov-chain model for a random walk so that the starting node is an *absorbing node*, i.e., a node from which the probability of a transition to any neighbour is 0 (and the probability of a transition to itself is 1). We then model the system as starting from the state distribution after the initial distribution, i.e., that in which all neighbours of the starting node have probability  $1/n$  where  $n$  is the number of neighbours per node.

**Definition 4** (*Markov-chain model with absorbing node*) Given a Markov-chain model as defined in Definition 3, let  $s < N$  be the position of the starting node. A Markov-chain model with absorbing starting node is defined in terms of a transition probability matrix  $M'$  and initial state probability distribution  $v^{(0)}$  as follows.

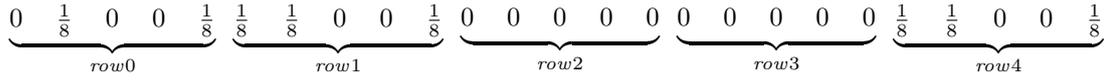


Figure 4: Row 0 of the matrix for a random walk on a network of  $5 \times 5$  nodes where *row0* to *row4* refer to the row in the network.

- The initial state probability distribution is that reached after 1 step from a distribution in which the agent is in the starting node with probability 1.

$$v^{(0)} = v M$$

where  $v_s = 1 \wedge (\forall j < N \bullet j \neq s \Rightarrow v_j = 0)$

- The starting node transitions to itself with probability 1 in  $M'$ .

$$M'_{s,s} = 1 \wedge (\forall j < N \bullet j \neq s \Rightarrow M'_{s,j} = 0)$$

- All other transitions in  $M'$  are as in  $M$ .

$$\forall i, j < N \bullet i \neq s \Rightarrow M'_{i,j} = M_{i,j} \quad \square$$

Due to the starting node being an absorbing node, the probability of being in the starting node never decreases as the number of steps increase. For example, for the 1-dimensional network above, the transition probability matrix is

$$M' = \begin{pmatrix} 0 & \frac{1}{2} & 0 & 0 & \frac{1}{2} \\ \frac{1}{2} & 0 & \frac{1}{2} & 0 & 0 \\ 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} \\ 0 & 0 & 0 & 1 & 0 \\ \frac{1}{2} & 0 & 0 & \frac{1}{2} & 0 \end{pmatrix}$$

and the state probability distribution is

$$\begin{aligned} v^{(0)} &= (0 \quad \frac{1}{2} \quad 0 \quad \frac{1}{2} \quad 0) \\ v^{(1)} &= (\frac{1}{4} \quad 0 \quad \frac{1}{2} \quad 0 \quad \frac{1}{4}) \\ v^{(2)} &= (\frac{1}{8} \quad \frac{1}{8} \quad \frac{1}{2} \quad \frac{1}{8} \quad \frac{1}{8}) \\ v^{(3)} &= (\frac{1}{8} \quad \frac{1}{16} \quad \frac{5}{8} \quad \frac{1}{16} \quad \frac{1}{8}) \\ v^{(4)} &= (\frac{3}{32} \quad \frac{1}{16} \quad \frac{11}{16} \quad \frac{1}{16} \quad \frac{3}{32}) \end{aligned}$$

and so on.

The probability of the starting node at step  $k$  in this model is the probability that the system has returned to the starting node within  $k$  steps. It can be used to calculate the coverage time as follows.

Let  $\gamma_i$  denote the expected number of new nodes covered in the  $i$ th step. The total expected number of nodes covered at the  $k$ th step,  $C_k$ , is then

$$C_k = \sum_{i=0}^k \gamma_i.$$

The initial node is covered at step 0, so we have  $\gamma_0 = 1$ . For all  $k > 0$ ,  $\gamma_k$  is equal to the probability that the node,  $n_k$ , reached at the  $k$ th step has not been visited before, i.e.,  $\gamma_k = P(n_k \notin \{n_i \mid i < k\})$  or

$$\gamma_k = P(n_k \neq n_0 \wedge \dots \wedge n_k \neq n_{k-1}).$$

Due to the regularity of the network and the fact that an agent behaves the same at each node, the probability of returning to the starting node after, say, 10

steps is equal to the probability of returning to the second node reached after 12 steps. More generally, we have  $P(n_{k-i} = n_0) = P(n_k = n_i)$ . From which it follows that  $P(n_{k-i} \neq n_0) = P(n_k \neq n_i)$ . Hence, from above

$$\gamma_k = P(n_0 \neq n_1 \wedge \dots \wedge n_0 \neq n_k)$$

which is equal to the probability that the system has not returned to the starting node,  $n_0$ , within  $k$  steps. In other words, given the modified Markov-chain model of Definition 4

$$C_k = \sum_{i=0}^k (1 - v_s^{(i)}).$$

Coverage time can then be defined as follows.

**Definition 5 (Coverage time)** Given a Markov-chain model with absorbing node as defined in Definition 4, the time to cover  $x\%$  of the nodes is the smallest  $k \geq 0$  such that

$$C_k = \sum_{i=0}^k (1 - v_s^{(i)}) = (x/100) * N. \quad \square$$

With directional bias, we add an additional dimension to our representation of a network: the current direction of movement. For a network with  $N$  nodes, the number of entries in the state probability distribution  $v$  is hence no longer  $N$  but  $n * N$  where  $n$  is the number of neighbours per node (and hence the number of directions of movement). Each entry represents the probability of being in a node having entered from a specific direction. We organise these entries so that those from  $d * N .. d * N + N - 1$  for  $0 \leq d < n$  denote the probabilities of being in a node of the network having entered from direction  $d$ . The corresponding transition probability matrix  $M'$  for the Markov-chain model is of size  $n * N \times n * N$ , and since there are  $n$  positions corresponding to the starting node (one for each direction from which the starting node was entered) there will be  $n$  absorbing positions in the matrix.

**Definition 6 (Markov-chain model for directional bias)** Let  $M$  be a transition probability matrix of size  $n * N \times n * N$  (such that all rows sum to 1). A Markov-chain model for directional bias is defined in terms of a transition probability matrix  $M'$  and initial state probability distribution  $v^{(0)}$  as follows.

- The initial state probability distribution is that reached after 1 step from a distribution in which the agent is in the starting node (with a particular current direction) with probability 1.

$$\begin{aligned} v^{(0)} &= v M \\ \text{where} \\ \exists d < n \bullet \\ &v_{d * N + s} = 1 \wedge \\ &(\forall j < n * N \bullet j \neq d * N + s \Rightarrow v_j = 0) \end{aligned}$$



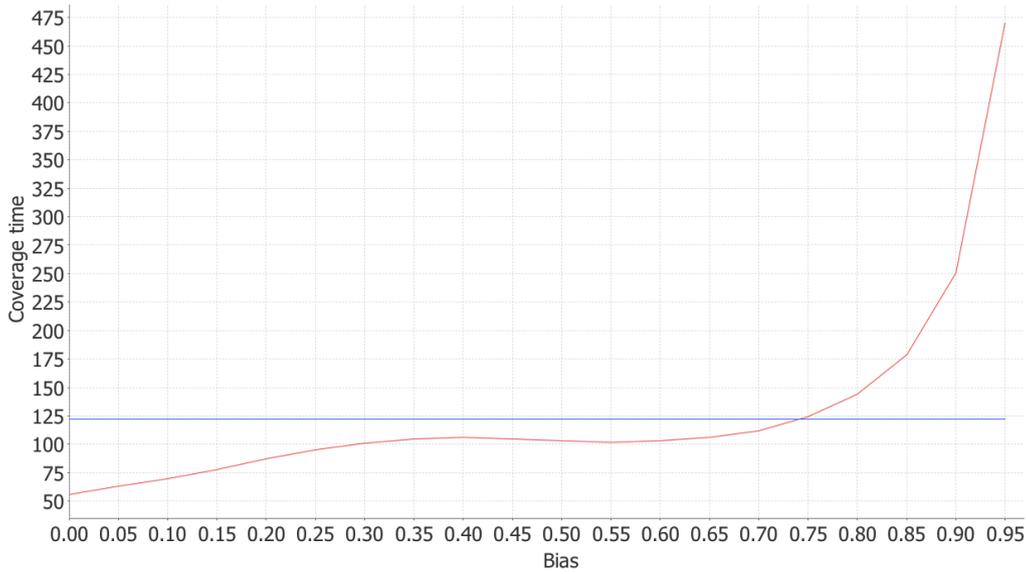


Figure 5: Random vs directionally biased walk for a 2-dimensional network of  $5 \times 5$  nodes.

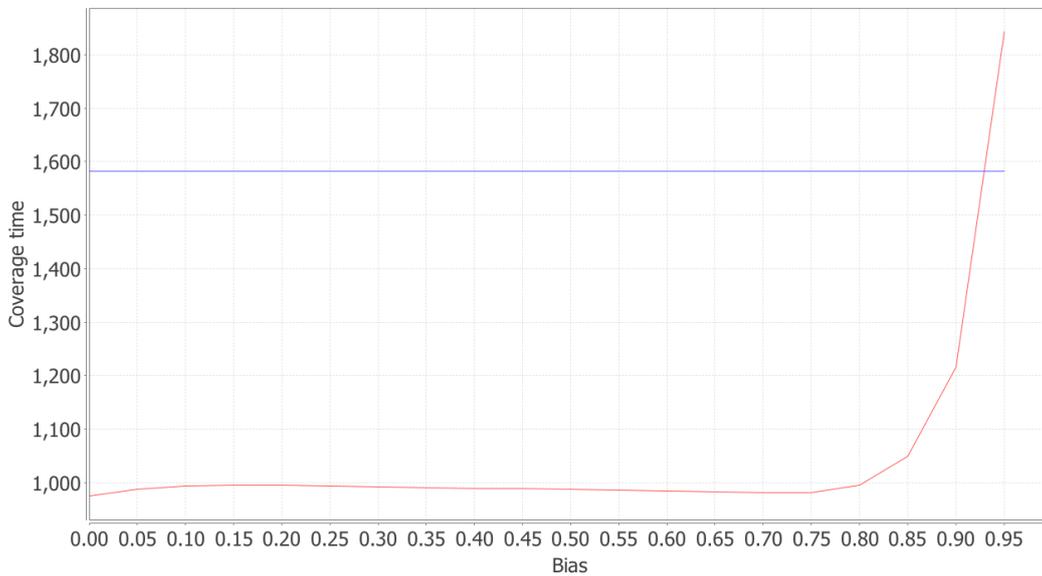


Figure 6: Random vs directionally biased walk for a 2-dimensional network of  $15 \times 15$  nodes.

micro-level simulation of a 2-dimensional network of  $100 \times 100$  nodes shows that a directionally biased walk (approximating a von Mises distribution with high  $\kappa$ ) is more efficient than a random walk.

The second part of our investigation considered the movement model of Definition 2, where occasional random steps are added to a directionally biased walk. Example plots for a network of  $5 \times 5$  nodes and the value of  $r$  set to 0.1 (an average of one random step in every 10) and 0.2 (an average of one random step in every 5) are shown in Fig. 7 and Fig. 8, respectively. Similar plots for a  $15 \times 15$  network are shown in Fig. 9 and Fig. 10.

The following results emerge from this analysis.

1. As may have been predicted, the addition of random steps moves the coverage time closer to that of a random walk. Hence, for bias values lower than the cross-over bias the coverage time increases, but for values higher than the cross-over value the coverage time decreases. Comparing Fig. 8 with Fig. 5 it can be seen that for a bias of 0 the coverage time has increased from around 55

to 70, and for a bias of 0.95 it has decreased significantly from around 470 to about 158. Comparing Fig. 10 with Fig. 6 it can be seen that for a bias of 0 the coverage time has increased from around 970 to 1050, and for a bias of 0.95 decreased significantly from over 1800 to around 1100.

2. The introduction of random steps increases the cross-over bias. For the  $5 \times 5$  network, with  $r = 0.1$  the cross-over bias increases to 0.78 (from 0.74 for no random steps) and for  $r = 0.2$  to 0.84. For the  $15 \times 15$  network, the cross-over bias increases from around 0.93 to a value higher than 0.95 (the last bias plotted) for both values of  $r$ .

## 5 Conclusion

In this paper we have investigated the effect of directional bias on the coverage time of random walks on regular, connected networks. Our analysis has shown that directional bias can reduce coverage time signif-

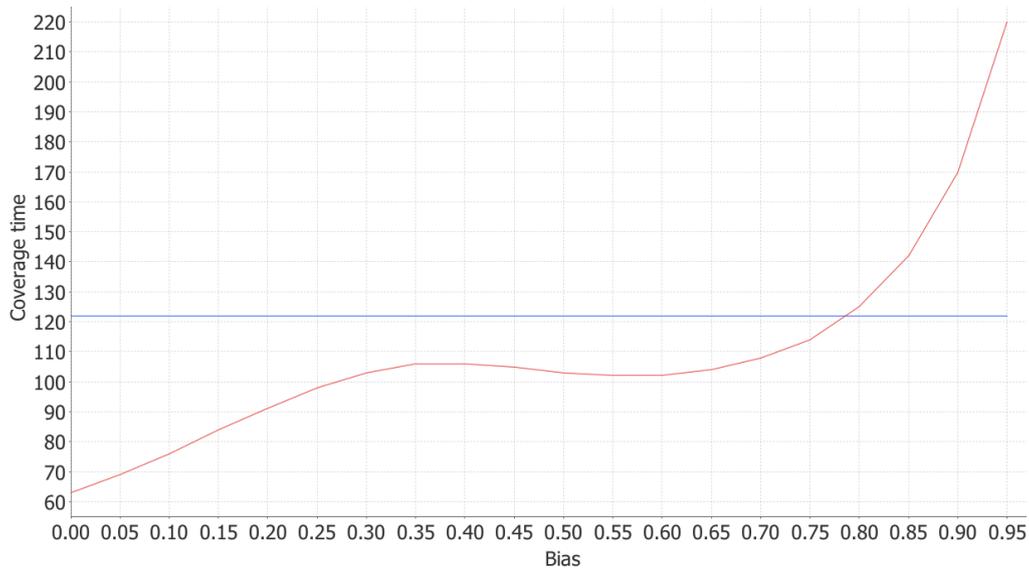


Figure 7: Random vs directionally biased walk with probability 0.1 of a random step for a 2-dimensional network of  $5 \times 5$  nodes.

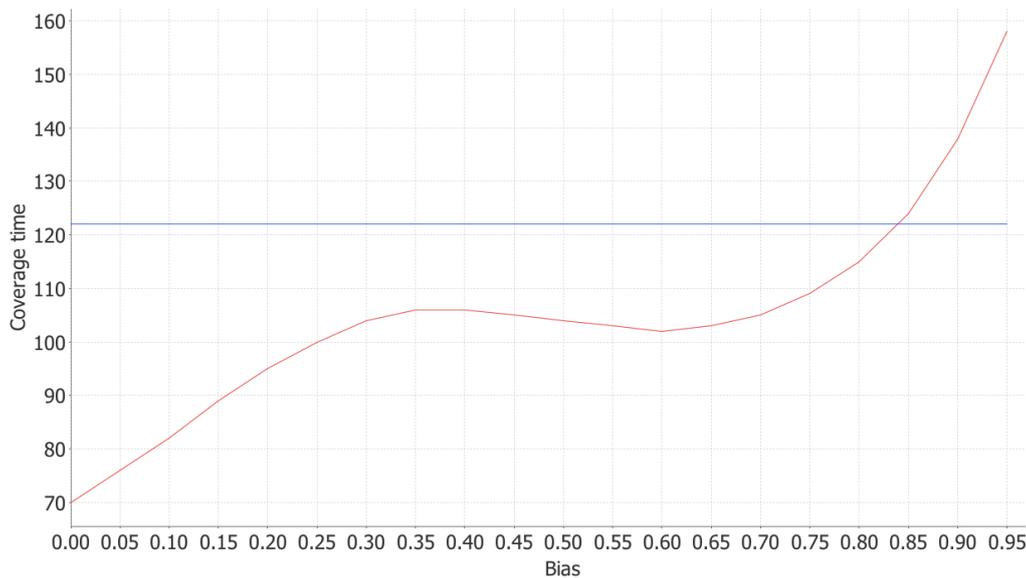


Figure 8: Random vs directionally biased walk with probability 0.2 of a random step for a 2-dimensional network of  $5 \times 5$  nodes.

icantly and has a greater effect the larger the network. However, this reduction occurs only when the bias (to continue in the same direction) is below a certain value we call the *cross-over bias*. The cross-over bias is dependent on the network size, increasing as the size of the network increases. Hence, high values of bias which work well at reducing coverage time in large networks, may be less effective, or even increase the coverage time, in smaller networks.

The cross-over bias can be increased by adding occasional random steps to a directionally biased walk — the more random steps, the higher the cross-over bias. Adding such steps, however, moves the coverage time of a directionally biased walk closer to that of a random walk (increasing coverage time for biases below the cross-over bias).

Our analysis also revealed that a movement model in which an agent changes direction by  $\pi/4$  radians (in either direction with equal probability) on each step is more effective in reducing coverage time than a standard directionally biased model. Further inves-

tigation of this, and similar models, is warranted.

Our investigation was facilitated by a macro-level model of random and directionally biased walks in terms of Markov chains. This model allowed us to calculate coverage time directly, in contrast to other approaches where coverage time is calculated as the average result obtained from numerous runs of a micro-level simulation. Although introducing a margin of error due to the limitations of floating-point arithmetic, our more abstract model has provided a practical means for obtaining a deeper, more complete analysis of directional bias than in previous work.

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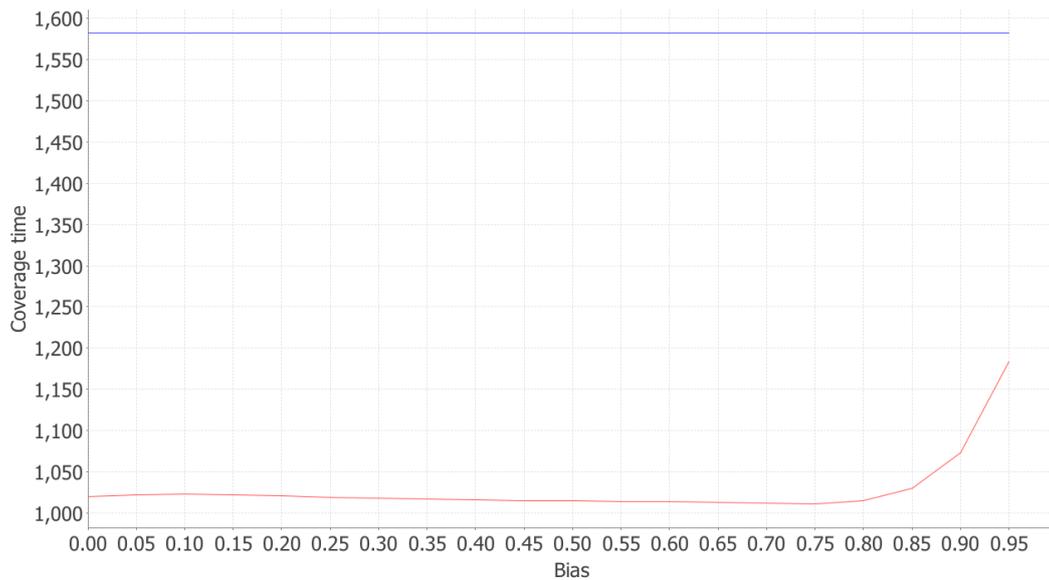


Figure 9: Random vs directionally biased walk with probability 0.1 of a random step for a 2-dimensional network of  $15 \times 15$  nodes.

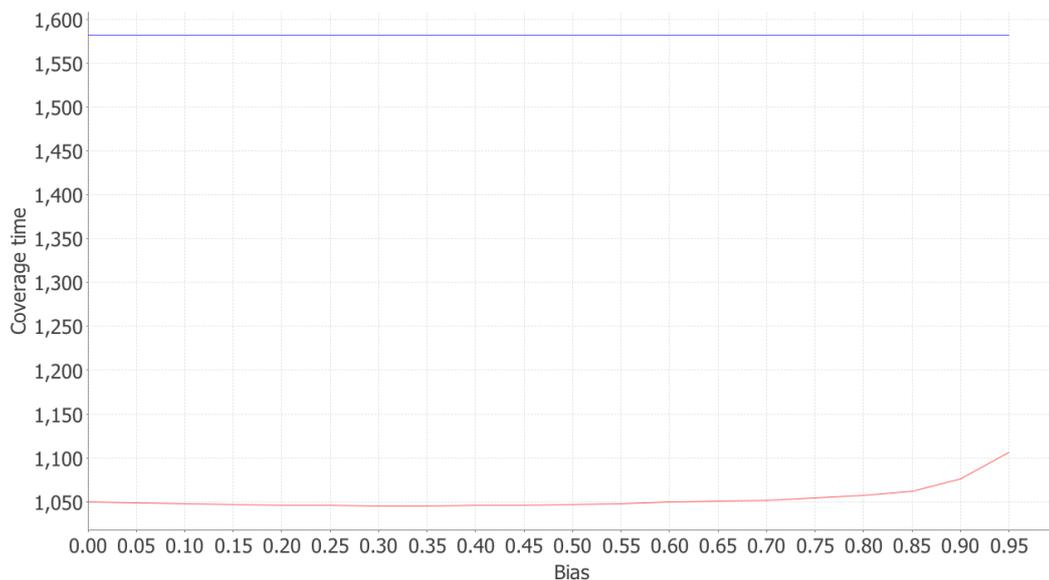


Figure 10: Random vs directionally biased walk with probability 0.2 of a random step for a 2-dimensional network of  $15 \times 15$  nodes.

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